

DIALOG

Elsinore 2023

Tacit Techniques

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These Tacit Techniques Will Blow Your Mind!

APL Wiki

apl.wiki/tacit#Tutorials

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xpqz.github.io/cultivations/Trains

Rich Park

youtu.be/Enlh5qwwDuY “Train Spotting”



What is the essence of tacit programming? Find out here!

Explicit code mentions arguments:

- Expression $(\uparrow / N) - \lfloor / N \leftarrow 3 \quad 1 \quad 4 \quad 1 \quad 5$
- Tradfn $\nabla R \leftarrow \text{Range } Y \dots$
- Dfn $\{ (\uparrow / \omega) - (\lfloor / \omega) \}$

Tacit code implies arguments:

- Tacit $\uparrow / - \lfloor /$



Shocking Revelation: You've been unknowingly using tacit constructs all along!

$f/$ f'' $\circ.g$ $f \setminus$ $A \circ g$ $f \boxplus$ $f \boxtimes B$

APL old-timers don't want you to know this one fact:

Operator application is actually just tacit programming!



Function composition

$f \circ g$ $f \cdot g$ $f \sim$ fgh $f \circ g$

Discover the mind-blowing secret:

Function composition is actually just plumbing!



Experts rave about tacit!

- Arguments in operands
- Memorable (like $\neq \subseteq \vdash$ and $+ \neq \div \neq$)
- Adjacency (like $\times -$ and $\vee / \underline{\epsilon}$)
- Brevity (like $F \ddot{\square} C$)
- DRY (Don't Repeat Yourself; like $\equiv \ddot{\rho}$)
- Just a general feeling of superiority and awesomeness

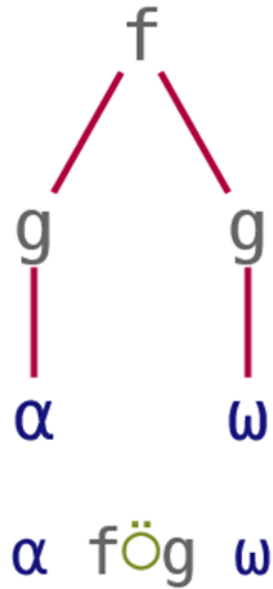


Over

The shape of an outer product $\alpha \circ f \omega$ is
 $(\rho\alpha) , (\rho\omega)$

We can write this as
 $\alpha , \ddot{\rho} \omega$

“pre-process both”



Beside

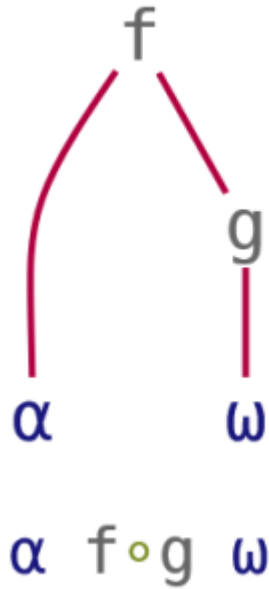
Location of α^{th} 1 in each element of ω is

$$\alpha \supset \underline{\tau} \omega$$

We can write this as

$$\alpha \supset \circ \underline{\tau} \omega$$

“pre-process right”

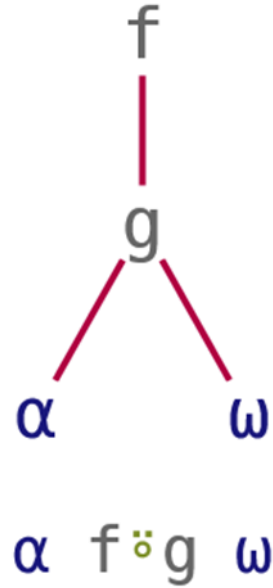


Atop

Any-presence of α in ω is
 $\forall \alpha \in \omega$

We can write this as
 $\alpha \forall \circ \in \omega$

“post-processing result”



Commute

A multiplication table of N is
 $(\iota\omega) \circ \cdot \times (\iota\omega)$

We can write this as
 $\circ \cdot \times \ddot{\sim} \iota\omega$

“selfie”



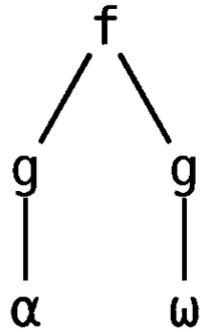
$f \ddot{\sim} \omega$



Tacit Techniques

Function Compositions

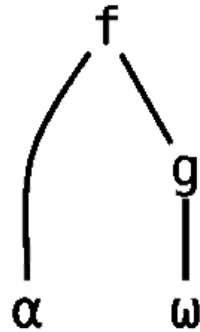
Over



$\alpha \text{ f } \ddot{\circ} \text{ g } \omega$

*pre-process
both*

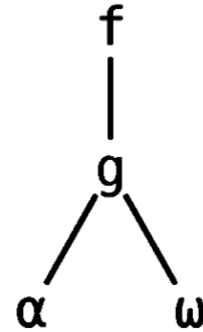
Beside



$\alpha \text{ f } \circ \text{ g } \omega$

*pre-process
right*

Atop



$\alpha \text{ f } \circ \text{ g } \omega$

*post-process
result*

Commute



$\text{f} \sim \omega$

selfie

Tasks: Tacify!

1. 'Hello' $\{(\square C \alpha) \equiv (\square C \omega)\}$ 'HELLO' 1
2. 'ab' 'cd' $\{\wedge / \alpha \in \omega\}$ " 'abba' 'dad' 1 0
3. 64 1000 0 $\{\alpha * \div \omega\}$ 3 4 10 0
4. 10 4 1 0 $\{\alpha \times \omega\}$ $^{-3}$ 2 0 $^{-1}$ $^{-10}$ 4 0 0



Train Introduction

Also a type of composition

Sequence of functions in isolation:

- Parenthesised:

$(+ \neq \div \neq) \ 3 \ 1 \ 4 \ 1 \ 5$

- Assigned

$\text{Avg} \leftarrow + \neq \div \neq$

$\text{Avg} \ 3 \ 1 \ 4 \ 1 \ 5$



Discussion: Writing Trains

(f Y) + (h Y)



Discussion: Writing Trains

$$(f \ Y) + (h \ Y) \rightarrow (f + h) \ Y$$

$$(f \ Y) + (\ Y)$$



Discussion: Writing Trains

$(f Y) + (h Y) \rightarrow (f + h) Y$

$(f Y) + (\quad Y) \rightarrow (f + \quad) Y$

$(X f Y) + (X h Y)$



Discussion: Writing Trains

$$(f Y) + (h Y) \rightarrow (f + h) Y$$

$$(f Y) + (\quad Y) \rightarrow (f + \quad) Y$$

$$(X f Y) + (X h Y) \rightarrow X (f + h) Y$$

$$(X \quad) + (X h Y)$$



Discussion: Writing Trains

$$(f Y) + (h Y) \rightarrow (f + h) Y$$

$$(f Y) + (\quad Y) \rightarrow (f + \quad) Y$$

$$(X f Y) + (X h Y) \rightarrow X (f + h) Y$$

$$(X \quad) + (X h Y) \rightarrow X (\quad + h) Y$$



Discussion: Writing Trains

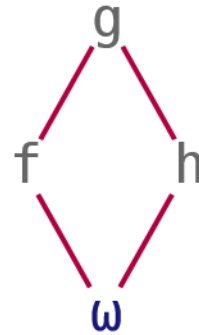
$(f Y) g (h Y) \rightarrow (f g h) Y$

$(f Y) g (\quad Y) \rightarrow (f g \vdash) Y$

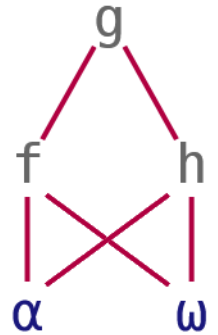
$g (h Y)$

$(X f Y) g (X h Y) \rightarrow X (f g h) Y$

$(X \quad) g (X h Y) \rightarrow X (\vdash g h) Y$



$(fgh)w$



$\alpha(fgh)w$



Discussion: Writing Trains

$$(f Y) g (h Y) \rightarrow (f g h) Y$$
$$(f Y) g (\quad Y) \rightarrow (f g \vdash) Y$$
$$g (h Y) \rightarrow (\quad g h) Y$$
$$(X f Y) g (X h Y) \rightarrow X (f g h) Y$$
$$(X \quad) g (X h Y) \rightarrow X (\vdash g h) Y$$
$$g (X h Y)$$


Discussion: Writing Trains

$(f Y) g (h Y) \rightarrow (f g h) Y$

$(f Y) g (\quad Y) \rightarrow (f g \vdash) Y$

$g (h Y) \rightarrow (\quad g h) Y$

$(X f Y) g (X h Y) \rightarrow X (f g h) Y$

$(X \quad) g (X h Y) \rightarrow X (\vdash g h) Y$

$g (X h Y) \rightarrow X (\quad g h) Y$



Discussion: Writing Trains

$$(f Y) g (h Y) \rightarrow (f g h) Y$$

$$(f Y) g (\quad Y) \rightarrow (f g \vdash) Y$$

$$f (g Y) \rightarrow (\quad f g) Y$$

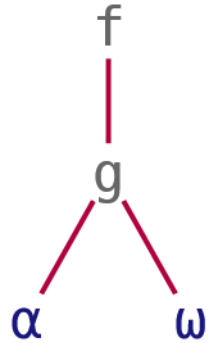
$$(X f Y) g (X h Y) \rightarrow X (f g h) Y$$

$$(X \quad) g (X h Y) \rightarrow X (\vdash g h) Y$$

$$f (X g Y) \rightarrow X (\quad f g) Y$$



$(fg)\omega$



$\alpha(fg)\omega$



Train Details

Some parts can actually be arrays:

- $A g h$ is $\{A\} g h$

Two sub-types of trains:

- Fork: $f g h$ and $A g h$
- Atop: $f g$

Longer trains:

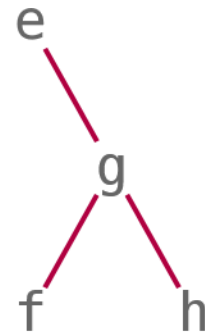
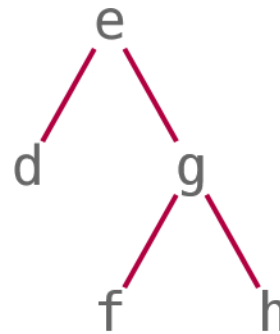
- $(d e f g h)$ is $(d e (f g h))$
- $(e f g h)$ is $(e (f g h))$



$(Agh)\omega$



$\alpha(Agh)\omega$



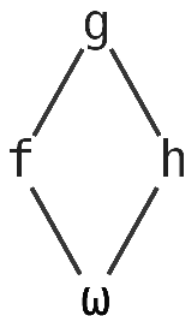
Tacit Techniques Trains

Monadic

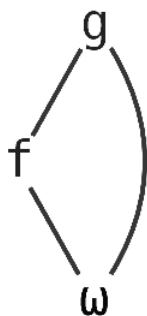
$$(f\ Y)\ g\ (h\ Y) \Leftrightarrow (f\ g\ h)\ Y$$

$$(f\ Y)\ g\ (\ \ Y) \Leftrightarrow (f\ g\ _)\ Y$$

$$g\ (h\ Y) \Leftrightarrow (\ \ g\ h)\ Y$$



$(fgh)w$



$(fgt)w$



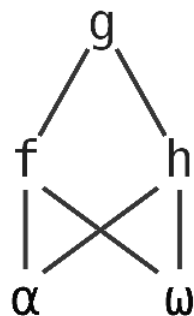
$(fg)w$

Dyadic

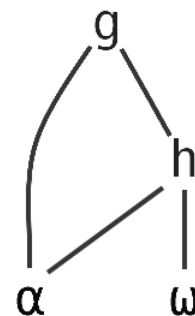
$$(X\ f\ Y)\ g\ (X\ h\ Y) \Leftrightarrow X\ (f\ g\ h)\ Y$$

$$(X\ \ \)\ g\ (X\ h\ Y) \Leftrightarrow X\ (_ \ g\ h)\ Y$$

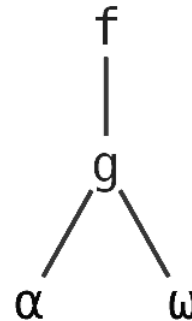
$$g\ (X\ h\ Y) \Leftrightarrow X\ (\ \ g\ h)\ Y$$



$\alpha(fgh)w$



$\alpha(-gh)w$



$\alpha(fg)w$

Example: Writing a Train

1 $\{ \wedge / (\Gamma \setminus \omega) = \omega \}$ 1 3 5 6 7

1 $\{ \wedge / (\Gamma \setminus \omega) = (\vdash \omega) \}$ 1 3 5 6 7

1 $\{ \wedge / (\Gamma \setminus = \vdash) \omega \}$ 1 3 5 6 7

1 $\{ (\wedge / \Gamma \setminus = \vdash) \omega \}$ 1 3 5 6 7

1 $(\wedge / \Gamma \setminus = \vdash)$ 1 3 5 6 7



Tasks: Tacify!

1. $\{2 \times \omega\}$ 2 7 1 8

4 14 2 16

2. 3 1 4 $\{(\alpha \cup \omega) \sim (\alpha \cap \omega)\}$ 1 6 1

3 4 6

3. $\{\cup \omega \vee \cap \omega\}$ 10

1 2 5 10

4. 2 $\{(\alpha \supset \Psi \omega) \supset \omega\}$ 3 1 4 1 5

4

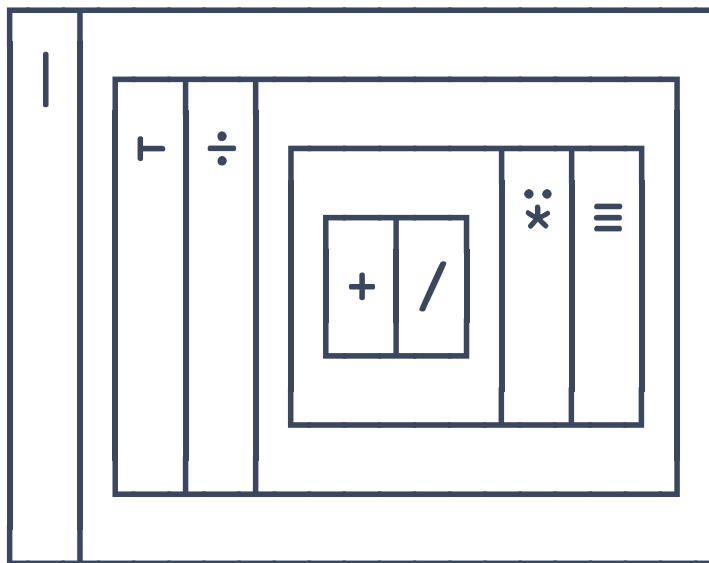


Amazing tool:]box on



]box on

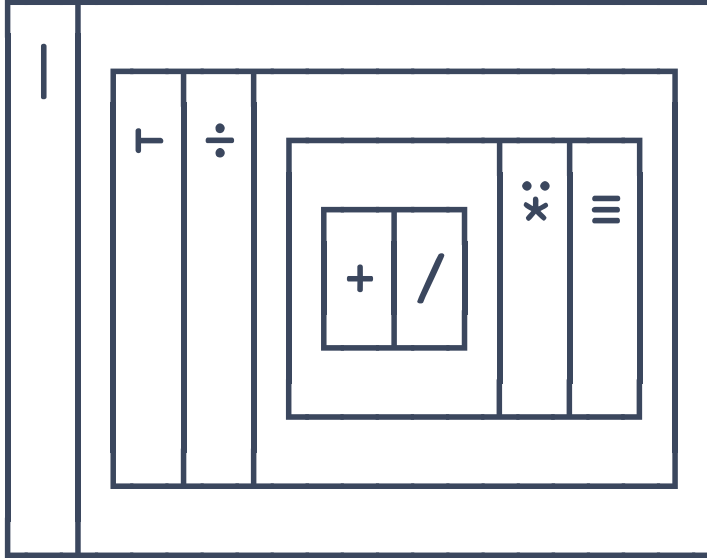
|┌÷+ / *≡



]box on -t=...

|┌÷+ / *≡

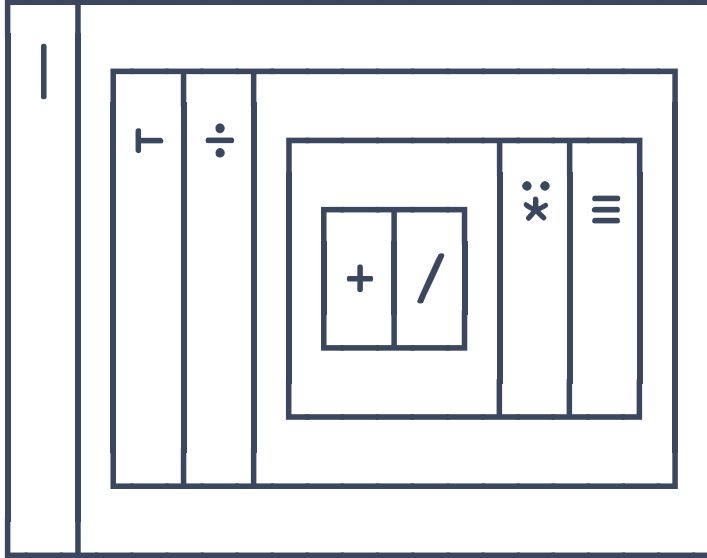
-t=box



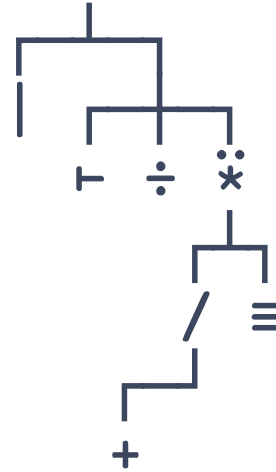
]box on -t=...

|┌÷+ / *≡

-t=box



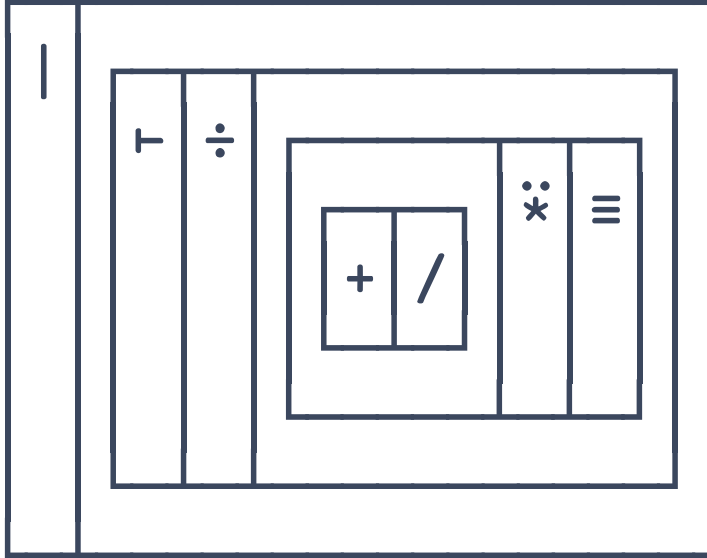
-t=tree



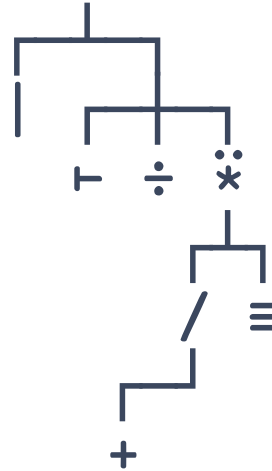
]box on -t=...

|┌÷+ / *≡

-t=box



-t=tree



-t=parens

| (┌÷((+ /) *≡))



Tasks: Tacify!

1. $\{(\theta\omega) \perp (\theta\omega)\}$ 1 1 1 0 1 1 0 0 3
2. ', ;' $\{(\sim\omega \in \alpha) \subseteq \omega\}$ 'ab, de; fgh' ab de fgh
nums ← 3 1 4 1 5
3. 4 $\{(\alpha + \neq \omega) \div \alpha\}$ nums 2.25 2.75
4. $\{(+ \neq \omega) \div \neq \omega\}$ nums 2.8
5. **Bonus task:** Combine 3 & 4 into an ambivalent function.



4 aspects of tacit that could ruin your life – #4 will blow your mind!

- ◆ Arguments in operands
- ◆ Monadic functions
- ◆ Recursion, Assignment, Dotting
- ◆ Selection (but 20.0...)



The ultimate paradox revealed: Arguments in operands

10 {x←α ◊ φ@{x<ω}ω} 1 2 13 14 5 16 7 18
 1 2 18 16 5 14 7 13

10 { φ@(α<⊢)ω} 1 2 13 14 5 16 7 18
 1 2 18 16 5 14 7 13

2 {(φ◊φ*α)ω} 3 3ρι9
 9 8 7
 6 5 4
 3 2 1



Inappropriate tacit discovered: Lots of monadic functions

$$\{\phi + \neq \uparrow \phi'' \omega\}$$

$$\phi \ddot{\circ} (+ \neq) \ddot{\circ} \uparrow \ddot{\circ} (\phi'')$$

$$\phi (+ \neq (\uparrow \phi''))$$

$$\phi + \neq \ddot{\circ} \uparrow \ddot{\circ} (\phi'')$$

$$\phi (+ \neq \ddot{\circ} \uparrow \phi'')$$



3 things tacit code just cannot do

Assignment

Namespace “dotting”

Recursion

$\{n; \phi \uparrow \alpha \pm \cdot n \leftarrow \alpha \cdot \square NL^{-2}\}$

$\{1 \geq \omega : 1 \diamond (\omega - 1) + \ddot{\nabla} (\omega - 2)\}$



The ugly truth about selection, and what we plan to do about it

$\{(3 > \omega) \neq \omega\}$ 3 1 4 1 5

4 5

$(3 \circ > \neq \vdash)$ 3 1 4 1 5

SYNTAX ERROR

$(3 \circ > \neq \vdash)$ 3 1 4 1 5

SYNTAX ERROR

$(3 \circ > \neq \ddot{\circ} \vdash)$ 3 1 4 1 5

4 5

$3 \circ > \underline{\circ} \neq$ 3 1 4 1 5



The ugly truth about selection, and what we plan to do about it

'aeiou' { $\omega[\alpha\Delta\omega]$ } 'hello world'
eohll wrld

'aeiou' { $\omega[\tilde{\alpha}\Delta\omega]$ } 'hello world'

LENGTH ERROR

'aeiou' { $\omega[\tilde{c}\alpha\Delta\omega]$ } 'hello world'
eohll wrld

'aeiou' ($c\ddot{o}\Delta[\text{r}]$) 'hello world'
eohll wrld

'aeiou' ($\Delta\geq\text{r}$) 'hello world'



Tasks: Convert tacit to dfn

1. Monadic $\times \times \lfloor \ddot{\circ} \rfloor$
2. Dyadic $\lfloor \circ \neq \uparrow \vdash$
3. Monadic $\equiv \ddot{\circ} (\square C \sim \circ ' ')$
4. Monadic $+ \neq \vdash > + \neq \div \neq$

Bonus tasks:

5. Monadic $\phi \equiv \vdash \vdash \equiv \phi$



Amazing tool: tacit.help

Transform tacit APL into dfn form

$f \leftarrow$

$f \ Y \Leftrightarrow$

$X \ f \ Y \Leftrightarrow$

Arrays: A, B, C, \dots Functions: a, b, c, \dots

+

#



7 hash tables that won't go away

$P \uparrow S$

$P \downarrow S$

$P \wr S$

$P \cup S$

$S \in P$

$S \sim P$

$S \cap P$

$P \circ \uparrow S$

$P \circ \downarrow S$

$P \circ \wr S$

$P \circ \cup S$

$(\epsilon \circ P) S$

$(\sim \circ P) S$

$(\cap \circ P) S$



7 hash tables that won't go away

```
s ← 'Hello, World!'
```

```
AVi ← □AV○ι ◇ {}AVi s
```

```
cmpx '□AVιs' 'AVi s'
```

```
□AVιs → 6.0E-7 | 0% □□□□□□□□□□□□□□□□□□□□□□□□
```

```
AVi s → 3.7E-7 | -39% □□□□□□□□□□□□□□
```

```
AVg ← □AV○ϕ ◇ {}AVg s
```

```
cmpx '□AVϕs' 'AVg s'
```

```
□AVϕs → 2.3E-6 | 0% □□□□□□□□□□□□□□□□□□□□□□□□
```

```
AVg s → 5.5E-7 | -77% □□□□□
```



Showdown: Memory vs CPU

$\bar{=}e$

$a (\wedge / \epsilon) b$
$a \{ \wedge / \alpha \epsilon \omega \} b$
$a \wedge / \epsilon b$
$a \{ \wedge / \alpha \epsilon \omega \} b$

$a \leftarrow (2e4p \text{ 'ab' 'cd' }) [?5e4p2]$
 $b \leftarrow \text{ 'aba' 'cad' } [?5e4p2]$

∇ $hwm \leftarrow mem \ expr$
 $\{ \} \square WA$
 $\{ \} 0(2000I)14$
 $hwm \leftarrow 2000I14$
 $\{ \} \perp expr$
 $hwm - \ddot{\sim} \leftarrow 2000I14$

∇



Showdown: Memory vs CPU

	($\bar{\phi}$, mem $\bar{\omega}$) e	
a	($\wedge/\bar{\epsilon}$) b	125835728
a	{ $\wedge/\bar{\alpha\epsilon\omega}$ } b	125835792
a	$\wedge/\bar{\epsilon}$ b	4194304
a	{ $\wedge/\bar{\alpha\epsilon\omega}$ } b	4194304
	($\bar{\phi}$, 1 \circ cmpx $\bar{\omega}$) e	
a	($\wedge/\bar{\epsilon}$) b	0.45675
a	{ $\wedge/\bar{\alpha\epsilon\omega}$ } b	0.4145
a	$\wedge/\bar{\epsilon}$ b	0.4245
a	{ $\wedge/\bar{\alpha\epsilon\omega}$ } b	0.42275



Showdown: Memory vs CPU

($\bar{\phi}$, mem) e

a (\wedge/ϵ) b 0

a { $\wedge/\alpha\epsilon\omega$ } b 0

a $\wedge/\circ\epsilon$ b 4194304

a { $\wedge/\alpha\epsilon\omega$ } b 4194304

($\bar{\phi}$, 1 \circ cmpx) e

a (\wedge/ϵ) b 0.0080625

a { $\wedge/\alpha\epsilon\omega$ } b 0.0078359375

a $\wedge/\circ\epsilon$ b 0.019703125

a { $\wedge/\alpha\epsilon\omega$ } b 0.021765625

$a \leftarrow 2 \uparrow a$



Check out this cool summary — you won't regret it!

Function composition:

- ◌ *Pre-process both*
- *Pre-process right*
- ◌◌ *Post-process*
- ◌◌◌ *Selfie*

Operators: long left scope

Trains: odd-even from right

Tools:]box on -t =...
tacit.help

Watch out for these:

Arguments in operands

Lots of monadic functions

Just don't try:

- ◆ Assignment
- ◆ Namespace “dotting”
- ◆ Recursion

Selection issues:

- ◆ Compress: $\vdash \ddot{\circ} \neq$
- ◆ Index: $\subset \ddot{\circ} \dots \square \dots$

