# How I won the APL Problem Solving Competition 

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## About Me

## Education

- Maths student at Università Degli Studi di Milano
- Abstract algebra
- Mathematical logic
- Computability
- wrote $\sim 50$ lines of Python at age 12 learned C in University programming language enthusiast $(\mathrm{C} / \mathrm{C}++$, Rust, Haskell, APL, BQN)


## My Journey With APL (and array programming)

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## First

MATLAB in university
Yes, it's technically an array language Interpreted, dynamic Weird syntax, this is nothing like C!

## Back to C (for now)

- Experiments with competitive programming
- Appreciation for arrays!


## My Journey With APL (and array programming)

Haskell from Tsoding's
community

- Composition patterns and combinators
- High information density

■
Similar to math notation


Language
Connor Hoekstra's Youtube channel

- Encountered APL
- Higher information density
- 

Terser math notation


## My Journey With APL (and array programming)

December 2022
learn APL by solving problems with it
$\sim 1$ new primitive per day
published my solutions on github and mastodon finished days $1-20$

## The Competition Problems

## The Competition Problems

Part 1 Bioinformatics
Task 5: Reading frame translation
"real world" programming
two very different solutions
Part 2 Potpourri

- Task 3: Time for a Change
- induction!
- opportunity for improvement


## Task 1.5: Reading frame translation

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Write a function that, given the name of a file in FASTA format, returns all the protein strings that can be translated from it, in all six reading frames.

## FASTA format

>Rosalind_2748
ATCAGGCTACCGTGTTTGCGGACGGGGGCTTAATCT CTTGTTGGCACAGCGGTGGCAGGAGGTCCCCGCCGA
'MVMATGVIVLNTRMRVTNDSNFGARYRGTCP' . . . ... 'MGL' 'MDRL' 'MRLPWSCLHIA'

## Final structure

$$
\text { orf }-\{c r f \ddot{a} \text { aasכфכreadFASTAW\} }
$$

- readFASTA performs the IO aas converts from DNA to list of reading frames crf extracts the protein strings


## Details

' \$ ' are stop codons.

- A protein string starts with ${ }^{\prime} \mathrm{M}^{\prime}$, ends before '\$'. Any 'M' not followed by '\$' doesn't start a protein.
' XXXMAAM\$YMBB' becomes:


## 'MAAM' 'M'

' XXXMAAM\$YMB\$' becomes:
'MAAM' 'M' 'MB'

## First Implementation of crf

Solve for one sequence, map and flatten.
(a -> [b]) -> [a] -> [b]

A utility: flat map modifier

$$
\_f \leftarrow\left\{\nu, / \alpha \alpha^{\prime \prime} \omega\right\}
$$

'abc'シ̈_f 25
'abcabcabcabcabc'

## First Implementation of crf

## Finding

## [Amino] -> [[Amino]]

$$
\begin{aligned}
& (' \$ ' \circ \neq \subseteq \vdash) \quad \text { 'AA\$BB\$XX' } \\
& \text { AA BB XX } \\
& \left((-' \$ ' \neq \vdash /) \downarrow^{\prime} \$ '^{\prime} \circ \neq \subseteq \vdash\right) \quad \text { 'AA\$BB\$XX' }
\end{aligned}
$$

$$
A A \quad B B
$$

Split each sequence by ' $\$$ '.
Drop the last split of sequences that don't end in '\$'.

## First Implementation of crf

## Suffixes

```
[Amino] -> [[Amino]]
```

＇M＇。（＝Сャ）＇XXMAAAMBB＇
MAAA MBBB
＇M＇。（，$\ddot{\sim} \backslash \ddot{\circ} \phi(=\subset \vdash)) ~ ' X X M A A A M B B '$
MBBB MAAAMBBB
c makes partitions starting with＇M＇．
Drops anything before the first＇ M ＇．
,$\ddot{\sim} \backslash \ddot{ } \phi$ concatenates the suffixes of the array．

## Putting it all together

## [[Amino]] -> [[Amino]]

$$
\begin{array}{ll}
\operatorname{crf} \leftarrow\{ \\
& p \leftarrow\left((-' \phi ' \neq \vdash /) \downarrow^{\prime} \phi^{\prime} \circ \neq \subseteq \vdash\right) \_f \omega \\
& U^{\prime} M^{\prime} \circ(, \ddot{\sim} \backslash \ddot{\circ} \phi(=\subset \vdash)) \_f p
\end{array}
$$

## Second Attempt at crf

## Index

## [[Amino]] -> ([Amino], [Int])

$$
s d-\left(, /,\left(\subset\left(+\backslash \neq{ }^{\prime \prime}\right)\right)\right), \circ ' \$ '^{\prime}{ }^{\prime} A \$ M B^{\prime} \quad \text { 'C\$MD\$' }
$$

$$
\text { A\$MB\$M\$MD\$\$ } 511
$$

## Second Attempt at crf

## Finding Begins

 and Ends$$
\begin{aligned}
& s=' A \$ M B \$ M \$ M D \$ \$ ' \\
& d=511
\end{aligned}
$$

$i \leftarrow \supset\{\uparrow \alpha(w[1+w \succeq \alpha])\} / ' M \$^{\prime}(\imath=){ }^{\prime \prime} \subset s$
( $\left.\begin{array}{lll}3 & 6 & 8\end{array}\right)\left(\begin{array}{llll}2 & 5 & 7 & 10\end{array}\right)$

## Second Attempt at crf

## Finding Begins

 and Ends$$
\begin{aligned}
& \mathrm{s}=\mathrm{I} A \$ \mathrm{MB} \$ \mathrm{M} \$ \mathrm{MD} \$ \${ }^{\prime} \\
& \mathrm{d}=511
\end{aligned}
$$

$$
\begin{aligned}
& \text { i-כ\{ta(w[1+wıa])\}/'M\$'(l=)"cs } \\
& \wedge \alpha \equiv 368 ; \omega \equiv 2571011
\end{aligned}
$$

$$
2
$$

$$
34
$$

## Second Attempt at crf

## Finding Begins

 and Ends$$
\begin{aligned}
& s=' A \$ M B \$ M \$ M D \${ }^{\prime}=1 \\
& d=511
\end{aligned}
$$

$$
\begin{aligned}
& \text { i-כ\{1a(w[1+wıa])\}/'M\$'(l=)"cs } \\
& 368 \\
& \begin{array}{lll}
5 & 7 & 10
\end{array}
\end{aligned}
$$

b e $\leftarrow^{-} 1+\downarrow(\sim i[2 ;] \in d) / i$
(5 7) (6 9)

## Second Attempt at crf

## Final Solution

$$
\begin{aligned}
& \text { crf }-\{ \\
& \text { s d-(,/, (c(+\尹"))), o'\$'" } \omega \\
& \text { i-כ\{ } 1 \alpha(\omega[1+\omega \underline{\imath} \alpha])\} / ' M \$ '(\underline{\imath}=) " \subset s \\
& \text { b e } \leftarrow^{-} 1+\downarrow(\sim i[2 ;] \in d) / i \\
& \text { ubl"et"cs } \\
& \text { \} }
\end{aligned}
$$

## Task 2.3: Time for a Change

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## Problem

Write a function that takes a list of denominations as a left argument and a total value as a right argument, and returns a matrix where each row represents a unique combination of the elements of the left argument that total the right argument.

In other words, find a non negative integer matrix $r$ with unique rows, maximising $\not \equiv r$ under the constraint $\omega \wedge .=r+. \times \alpha$

## Task 2.3: Time for a Change

Pruning

$$
\operatorname{GCD}\left(\alpha_{1}, \cdots, \alpha_{n}\right) \mid \omega \Leftrightarrow \exists \beta_{1}, \cdots, \beta_{n} \in \mathbb{Z}, \omega=\sum_{i=1}^{n} \alpha_{i} \beta_{i}
$$

## Inductive base

$$
\begin{array}{ll}
0 \neq \omega \mid \ddot{\sim} v / \alpha: & (0, \not \equiv \alpha) \rho 0 \\
1=\not \equiv \alpha: & ; \omega \div \supset \alpha
\end{array}
$$

- Avoid recursive calls

Base case is trivial

## Task 2.3: Time for a Change

## Inductive step

$$
\alpha \equiv 1 \quad 2 \quad 3 \diamond \omega \equiv 7
$$

$$
\begin{array}{lll} 
& i \leftarrow 0,2\lfloor\omega \div a \leftarrow \supset \phi \alpha \\
0 & 1 & 2
\end{array}
$$

## Task 2.3: Time for a Change

## Solving

## subproblems

$$
\alpha \equiv 123 \diamond \omega \equiv 7 \diamond i \equiv 012 \diamond(w-a \times i) \equiv 741
$$

|  | $s \leftarrow\left(c^{-} 1 \downarrow \alpha\right) \nabla " \omega-a \times i$ |  |  |  |  |
| :--- | :--- | :--- | :--- | :--- | :--- |
| 7 | 0 | 4 | 0 | 1 | 0 |
| 5 | 1 | 2 | 1 |  |  |
| 3 | 2 | 0 | 2 |  |  |
| 1 | 3 |  |  |  |  |

## Task 2.3: Time for a Change

## Merging step

$$
\left.\right)
$$

## Task 2.3: Time for a Change

## Final solution

$$
\begin{aligned}
& \text { makeChange } \leftarrow\{ \\
& 0 \neq \omega \mid \ddot{\sim} v / \alpha: \quad(0, \not \equiv \alpha) \rho 0 \\
& 1=\not \equiv \alpha: \quad ; \omega \div \supset \alpha \\
& i \leftarrow 0, \imath\lfloor\omega \div a \leftarrow د \phi \alpha \\
& s \leftarrow\left(c^{-} 1 \downarrow \alpha\right) \nabla{ }^{\prime \prime} w-a \times i \\
& \nu_{i} / s,{ }^{\prime} \text { i } \\
& \text { \} }
\end{aligned}
$$

Algorithmically better solution discussed on my blog

## What I leaned

## What I leaned

APL
is not as opinionated as I thought.

- direct code translation can work
- functional patterns apply

Array
Programming
has been there all along
APL techniques apply to MATLAB, $\mathrm{C}++$, etc.

- It changes the way you think

